

Research seminar week 8

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This week:

- Background in learning theory
- Niyogi, chapter 2 (3 and 4).

Framework for Learning

- Finite *alphabet* Σ .
- *Language* $L \subset \Sigma^*$.
- Grammar g , generating language L_g .
- Family of grammars \mathcal{L} ,
family of languages $\mathcal{L} = \{L_g | g \in \mathcal{G}\}$.

NB: Language not an input-output mapping now.

Framework for Learning

- *Example sentences* s_i : from target L_t .
- Set of possible example sentence sequences $\mathcal{D} = (\Sigma^*)^*$.
- *Hypothesis languages* $h \in \mathcal{H}$.
- Learning algorithm \mathcal{A} : effective procedure $\mathcal{A} : \mathcal{D} \rightarrow \mathcal{H}$.

Framework for Learning

- $d(\cdot, \cdot)$: distance of grammars/languages.
- Criterion of success:

$$\lim_{n \rightarrow \infty} d(g_t, h_n) = 0$$

where $h_n = \mathcal{A}(s_1, \dots, s_n)$.

Approaches to Learning

- Inductive inference / identification in the limit / Gold-learning.
- Probably Approximately Correct (PAC) learning

Identification in the limit

- *Text* t for language L : infinite sequence s_1, \dots, s_n, \dots , s.t. each $s_i \in L$, and all elements of L appears at least once in t .
- t_k : first k elements of text t .
- $t(k) = s_k$.

Identification in the limit

- Learning algorithm \mathcal{A} identifies (learns) target g_t on text t *in the limit*, if $\lim_{k \rightarrow \infty} d(\mathcal{A}(t_k), g_t) = 0$.
- *cal* \mathcal{A} identifies g_t in the limit, if it identifies g_t in the limit for all texts of L_{g_t} .
- Family \mathcal{G} is *identifiable in the limit* if there is an algorithm \mathcal{A} that identifies every $g \in \mathcal{G}$ in the limit.

Identification in the limit

- Gold's theorem (1967): family consisting of all finite languages and at least one infinite language is not learnable in the limit.
- Not learnable: regular languages; context free languages; infinite regular languages.
- Poverty of Stimulus; nativist arguments.

PAC Learning

Probably Approximately Correct Learning

(Vapnik and Chervonenkis 1971)

- “Probably”: on “almost every” sequences of data.
- “Approximately correct”: the algorithm gets close enough to target.

PAC Learning

Results:

- PAC unlearnable: all finite languages; regular languages; context free languages.

Complexity of learning

- Speed of convergence.